

# Session 8: IO and interaction

COMP2221: Functional programming

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- Discussed Haskell's implementation of expression evaluation: lazy evaluation
- Saw how lazy evaluation allows for programming with infinite data structures
- Discussed difference between *strict* and *lazy* evaluation, and how to implement strict functions in Haskell.

# IO and side effects

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# Batch programs

- So far, we've only written *batch* programs
- That is, programs that take all their inputs at the start and provide output at the end.

## Batch programs



- To change what we compute, need to change source code and rerun.
- What if want to write programs that allow *interaction*?

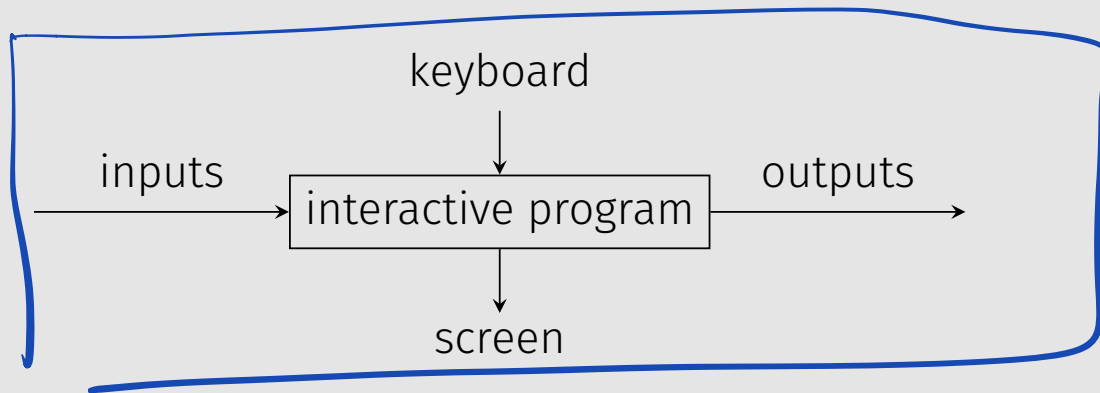
# Interactive programs

- What if we want to use Haskell to write interactive programs?
- These read from the “keyboard” and write to the “screen” as they are running

*“websocket” / video camera*

## Interactive program

*↳ this process pure?*



- Now the results can change depending on what input is provided (via the keyboard)  $\Rightarrow$  problem?

# A problem

- Haskell programs are *pure* mathematical functions
- ⇒ Haskell programs therefore *have no side effects*

## Definition (Side effect)

Modify some (internal/hidden) state as well as returning a value

- Reading from the keyboard and writing to the screen are side effects
- ⇒ Interactive programs *have side effects*
- How to square this circle?

Oh no!  $y = (\text{putChar 'x'}, \text{putChar 'x'})$   
 $y'z \text{ let } x = \text{putChar 'x'} \text{ in } (x, x)$

# Conceptual idea

- We can think of an interactive program as a pure function of type **World -> World**
  - That is, it takes the current *state of the world* as input and produces a *modified world* as output
- ⇒ new **World** object reflects any side effects that were performed

## IO actions

```
type IO a = World -> (a, World)
```

Input/output eats the world and produces a result of type `a`, along with a new world.

output:            main :: IO ()  
input              readString :: IO String

# A solution: *actions*

- Copying the world is too expensive in practice
- ⇒ Introduce *new types* to distinguish pure expressions from *impure actions*
- ⇒ Use the concept, but Haskell uses a *primitive type*: implementation details are hidden.
- These actions may have side effects
- Now we can write interactive programs in Haskell and “hide” the side effects behind a type.

## The **IO** type

```
data IO a = ... -- "Opaque" implementation
```

The type of actions that return a value of type **a**.



# Basic actions

## Reading

```
getChar :: IO Char
getChar = ...
```

Read a character from the keyboard, echo it to the screen, and return it

*Read and echo a character.*

*send output of getChar  
w/ putChar.*



*type mismatch.*

## Writing

```
putChar :: Char -> IO ()
putChar c = ...
```

Write a character to the screen and return nothing (indicated by the empty tuple)

# Bridging from expressions into actions

- For type safety, we need a way of “wrapping” values into actions
- Allows us to bring side-effect-free expressions into the “action” world.

## From pure to impure

```
return :: a -> IO a  
return v = ...
```

“Lift” a pure expression into an impure action.

Note: no way of turning an action back again.

can't unwrap  
(safely).

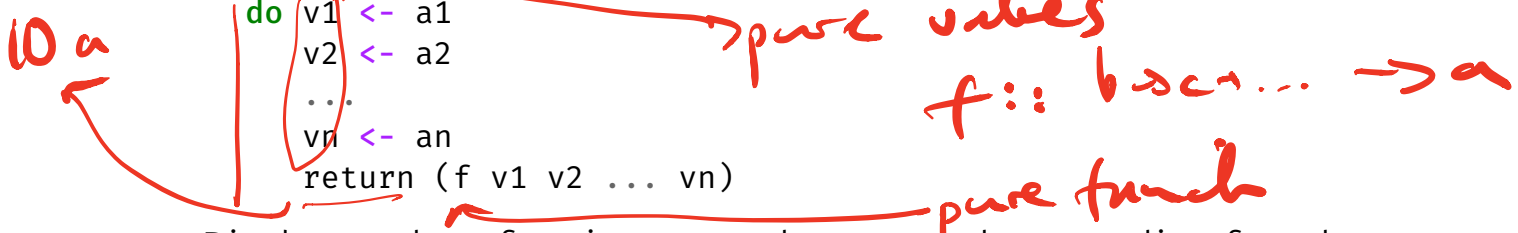
## WARNING!

The name `return` is rather misleading when coming from imperative languages. Calling `return` *does not* affect control flow.

# Sequencing actions

- We can combine a sequence of **IO** actions using **do** notation

```
do v1 <- a1
   v2 <- a2
   ...
   vn <- an
   return (f v1 v2 ... vn)
```



- Binds results of actions to values and then applies  $f$  to the values and lifts into “action-land” with `return`.

## Similarity with list comprehensions

- Each expression `vi <- ai` is called a *generator*
- If we want to execute an action, but don't care about the result, we can use `_ <- ai` or just `ai`

# Example: reading characters

## A first action

```
act :: IO (Char, Char)
act = do x <- getChar
        getChar
        y <- getChar
        return (x, y)
```

- Read three characters, discard the second, and return the first and third.
- Note use of `return`, without it we would get a type error  
 $\Rightarrow (x, y) :: (\text{Char}, \text{Char})$ , but we need an `IO (Char, Char)`.

# More primitives

## Read a string

```
getLine :: IO String
getLine = do x <- getChar
            if x == '\n' then
                return []
            else
                do xs <- getLine
                   return (x:xs)
```

# More primitives

## Write a string

```
putStr :: String -> IO ()
putStr []      = return ()
putStr (x:xs) = do putChar x
                   putStr xs
```

## Write a string with a new line

```
putStrLn :: String -> IO ()
putStrLn xs = do putStr xs
                 putChar '\n'
```

# When is an action performed?

- Actions never require arguments: `act :: IO a` is not a function
- Just specify that something will be done

⇒ Must be “run” to execute

- GHCi knows to run actions at the prompt

```
Prelude> x = putStrLn "hello"
```

*x :: IO ()*

```
Prelude> x
```

```
hello
```

```
Prelude> x
```

```
hello
```

- Conversely, when writing a program to be compiled, GHC only ends up running the main action.

⇒ Compare `main` function in C/Java.

*module Main where*

*main :: IO ()*

# Why these complications?

- One might wonder why we can't write actions as functions
- They would then behave like we're "used to"

## Why not this?

```
getChar :: () -> Char  
getChar _ = ...
```

“`getChar` ignores its argument and returns a `Char`”

- The problem is one of purity and *referential transparency*



# Pure vs. Impure

## Pure

- Always produces same result when applied to the same arguments
- Never has side effects
- Never alters state

## Impure

- *May* produce different results when applied to the same arguments
- *May* have side effects
- *May* alter state

- Impure functions are not *referentially transparent*

## Definition (Referential transparency)

Replacing an expression by its *value* does not change the behaviour of the program

- Not possible with `getChar`: which `Char` should we substitute?  
⇒ Can't treat them as normal (pure) functions

# Actions as promises

- To fix the issue of referential transparency, **IO** is introduced
  - We can think then of a type **IO Char** as a *placeholder* for a **Char** that will only materialise once the program executes
  - Moreover, it encapsulates a *promise* that this **Char** will actually appear.
- ⇒ manipulating an **IO Char** is equivalent to setting up “plans” to be executed when the **Char** materialises.
- This way, we maintain type safety “inside” the action.

IO Monad, oh it's easy, just  
think of the list monad.

# An example interactive program

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# Hangman

- Let's write a simple “hangman” game:
- Player A secretly enters a word
- Player B tries to figure out the word with a sequence of guesses
- For each guess, the program indicates which letters of the secret word are in the guess
- Game is over once the guess is correct
- Let's implement this “top down”

# Hangman I

- We start by importing useful IO functions
- The `main` function will just run the game

```
import System.IO
-- Set terminal output buffering so we see prints immediately
main = do hSetBuffering stdout NoBuffering
          hangman
```

- We prompt for a word, read it secretly (without echoing) and then run the play loop.

```
hangman :: IO ()
hangman = do putStrLn "Think of a word: "
            word <- secretlyGetLine
            play word
```

# Hangman II

- Now we want to read input from the terminal, but without echoing
- `getLine` does the former, but also echos as we type.
- Here we turn off the echoing and instead print hyphens

```
secretlyGetLine :: IO String
secretlyGetLine = do hSetEcho stdin False
                    xs <- getLine
                    putChar (replicate (length xs) '-')
                    hSetEcho stdin True
                    return xs
```

- Notice how *inside* the `do` block, the *results* of actions are just normal pure types.

# Hangman III

- Finally, we define how to play the game
- We repeatedly ask for a guess, either it was correct

```
play :: String -> IO ()
play word = do putStr "What is your guess? "
               guess <- getLine
               if guess == word then
                 do putStr "Correct! The word was "
                    putStrLn word
```

- Or it was not, in which case we show which letters matched and prompt again.

```
else
  do putStrLn (match word guess)
     play word
```

```
match :: String -> String -> String
match xs ys = [if x `elem` ys then x else '-' | x <- xs]
```

# Building block summary

- Prerequisites: none
- Content
  - Saw **IO** action, and how this allows side-effectful input and output in Haskell programs
  - Discussed difference between pure and impure functions
  - Saw *sequencing* and **do** syntax for **IO** actions
  - Saw how to write interactive programs that prompt for input from terminal.
- Expected learning outcomes
  - Student can *explain* how Haskell deals conceptually with side-effectful IO.
  - Student can *write* simple interactive programs
- Self-study
  - None