

# Session 2: Types and classes

COMP2221: Functional programming

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- What are some differences between functional and imperative programming?
- Which programming model more closely mirrors the way computers execute?
- What are interpreters and compilers? (in very broad terms)
- What are some advantages of an interpreter?
- Side effects (definition)
- Why can side effects easily introduce bugs?

# Types

On a computer: everything is bits. → types tell us how to interpret the bits.

- Mathematics and programming rely on the notion of types
- Tell us *how* to interpret a variable
- Provide restrictions on valid *operations*

## Example: Java/C

```
int a = 4;
```

```
int b = 3;
```

```
double c = a/b;
```

```
double a = 4;
```

```
double b = 3;
```

```
double c = a/b;
```

Integer division and floating point division use different instructions. → types tell us which instruction to use.

# Types

- Mathematics and programming rely on the notion of *types*
- Tell us *how* to interpret a variable
- Provide restrictions on valid *operations*

## Example: Java/C

```
int a = 4;
```

```
int b = 3;
```

```
double c = a/b;
```

```
double a = 4;
```

```
double b = 3;
```

```
double c = a/b;
```

Result depends on input types.

Since computers represent *everything* as sequences of bits, types are also required to define what these bit streams mean.

...required to know what a bit sequence means.

## Implementation

Find the correct implementation of, for example, '+'

## Correctness

check whether an operation on some data is valid and/or well-defined.

check whether a code fragment is correct (type safety)

## Documentation

document the code's semantics (for the reader)

# Types in Haskell

## Haskell is

Strongly, statically typed.

⇒ every well-formed expression has exactly one type, these types are known at *compile time*

## Definition (Type)

A type identifies a *collection* of values

## Example

- **Bool** the two logical values **True** and **False**.
- **Bool**  $\rightarrow$  **Bool** the set of all functions that take a **Bool** as input and produce a **Bool** as output.
- We will see more standard types soon

# Notation and inspection

## Attaching types

Haskell's notation for “e is of type T” is spelt

```
e :: T
-- False is of type Bool
False :: Bool
-- not is of type Bool -> Bool
not :: Bool -> Bool
```

## What type does X have?

Every valid expression in Haskell must have a valid type.

You can ask GHCi what the type of an expression is with the command `:type expr`

```
Prelude> :type sum
sum :: Num a => [a] -> a
```

# Type checking I

- Translators must check for *type correctness*

## Definition (Statically typed language)

We check correctness at translation time. (C/Java/Haskell/...)

⇒ invalid types mean “translation error”

```
-- Invalid  
foo :: a -> Int  
foo f = 1 + f
```

*Incorrect programs are spotted early~*

## Definition (Dynamically typed language)

We check correctness at run time. (Python “duck typing”)

⇒ invalid types only detected if we “use them”

```
# Fine as long as f supports addition with a number  
def foo(f):  
    return 1 + f
```

*Have non-numeric return types.*



# Type checking II

- How does the translator determine the type of an expression?

## Explicit annotation

Programmer annotates all variables with type information  
(e.g. C/Java)

*C++ you can use auto  
But only local. →*

## Type inference

Translator *infers* the types of variables based on the operations used (e.g. Haskell/ML)

*Hindley-Milner-style type inference.*

## Duck typing

Translator/runtime just tries the operation, if it succeeds, that was a valid type! (Python)

*→ link to some commentary on this on the website.*

# Demo time

Let's look at some types

# Building block summary

- Prerequisites: none
- Content
  - Different concepts of typing (dynamic/static)
  - Looked at some builtin Haskell types
  - Looked at list and tuple types
- Expected learning outcomes
  - student *knows* names of basic Haskell types compiling a programming language
  - student can *explain* difference between lists and tuples in Haskell.
  - student can *use* the Haskell interpreter to determine the type of an expression.
- Self-study
  - None

Functions have types

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# Programming with functions

- Functions have types in all programming languages, Haskell makes this particularly explicit

## Functions of one argument “unary”

Map from one type to another

```
not  :: Bool -> Bool  
and  :: [Bool] -> Bool
```

## Functions of two arguments “binary”

Map from two types to another

```
add  :: (Int, Int) -> Int
```

“add eats two Ints and returns an Int”

# An alternative view

- Since functions are *first class objects*, functions of *more than one* argument are typically written in Haskell as *functionals*
- Naturally extends from binary to n-ary functions

## “Curried” view of binary functions

```
add :: Int -> (Int -> Int)
```

“add eats an Int and returns a function which eats an Int and returns an Int”

- This idea comes from the formalism of Lambda calculus

# Currying

## Definition (Currying (informal))

Turn a function of  $n$  arguments into a function of  $n - 1$  arguments.

## History

- Idea first introduced by Gottlob Frege
- Developed by Moses Schönfinkel in the context of combinatory logic
- Further extended by Haskell Brooks Curry working in logic and category theory
- Name “currying” coined by Christopher Strachey (1967)

## Why currying?

- *easier* to reason about and prove things with functions of only 1 variable!
- Flexibility in programming: makes composing functions simpler
- Related to *partial evaluation* where we bind some variables in an  $n$ -ary function to a value

# Demo time

Let's look at some functions



# Building block summary

- Prerequisites: none
- Content
  - Specifying input and output types of functions
  - Functions have types, and so returning functions is natural
  - Functions of multiple variables can be defined using tuples, or else returning functions on a reduced parameter list
  - Introduction to currying
- Expected learning outcomes
  - student *knows* how to specify the type of a function
  - student *knows* two ways of writing functions of multiple arguments.
  - student can *explain* the difference between these paradigms (currying)
  - student can *illustrate* where currying or not makes a difference in semantics of function application
- Self-study
  - Lecture code.